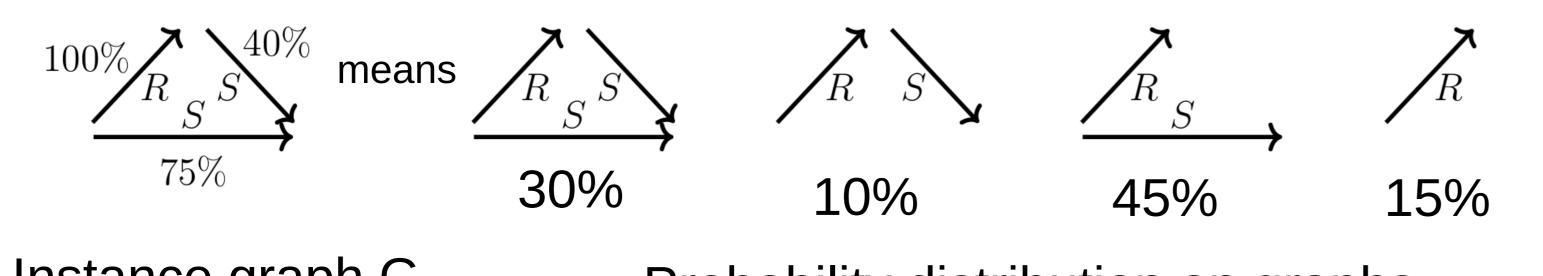
# **Probabilistic Graph Homomorphism: Combined Complexity**

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#### **Probabilistic Graph Homomorphism**

### **Probabilistic instance graph**:

- Each edge is **present** or **absent** with given probability
- Independence across edges



#### **Known Results about Data Complexity**

# **Data complexity**:

- Fix query graph G
- Study the complexity as a function of |G|

# Dalvi & Suciu [1] imply:

- There is a class S of **safe** query graphs
- → Data complexity is **PTIME** if  $G_{\cap} \in S$

Instance graph G Probability distribution on graphs

Query graph  $G_{\Omega}$  K

# **Probabilistic graph homomorphism:**

• **INPUT:** query graph  $G_o$  and probabilistic instance graph  $G_i$ • OUTPUT: probability that  $G_0$  has a homomorphism to  $G_1$ 

**Example:** for  $G_{O}$  and  $G_{I}$  above, prob. = 30% + 10% = 40%

→ Data complexity is **#P-hard** if  $G_{\cap} \notin S$ 

Amarilli & al [2] imply:

•  $G_k = all graphs of treewidth < k$ 

 $\rightarrow$  Data complexity of any query on  $G_{\nu}$  is linear-time

What about the **combined** complexity? (complexity as a function of both |G| and  $|G_0|$ )

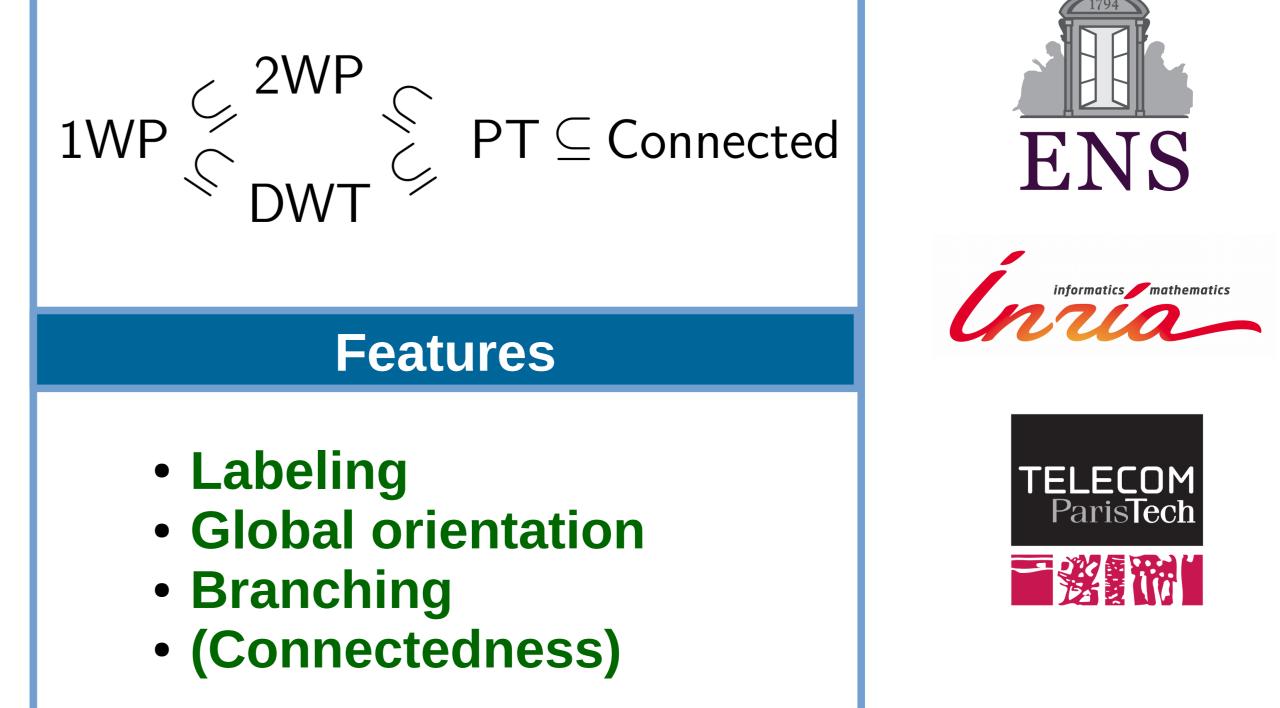
#### Problem

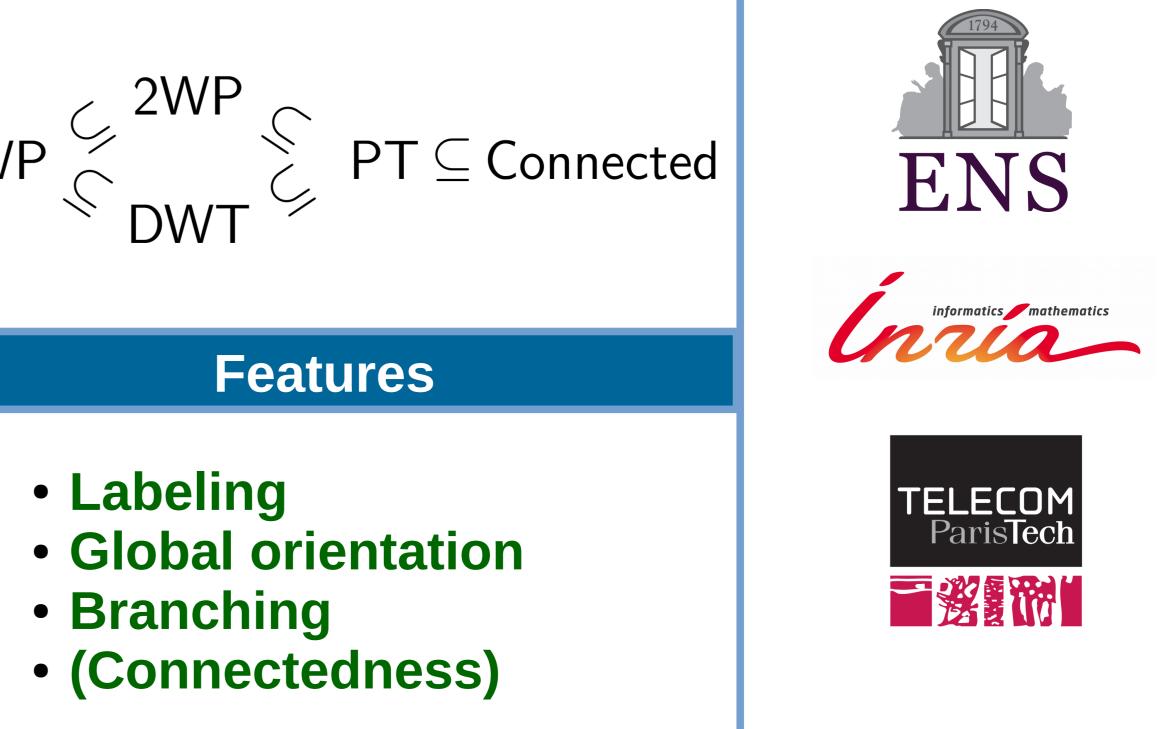
# Study the **combined** complexity of the probabilistic graph homomorphism problem

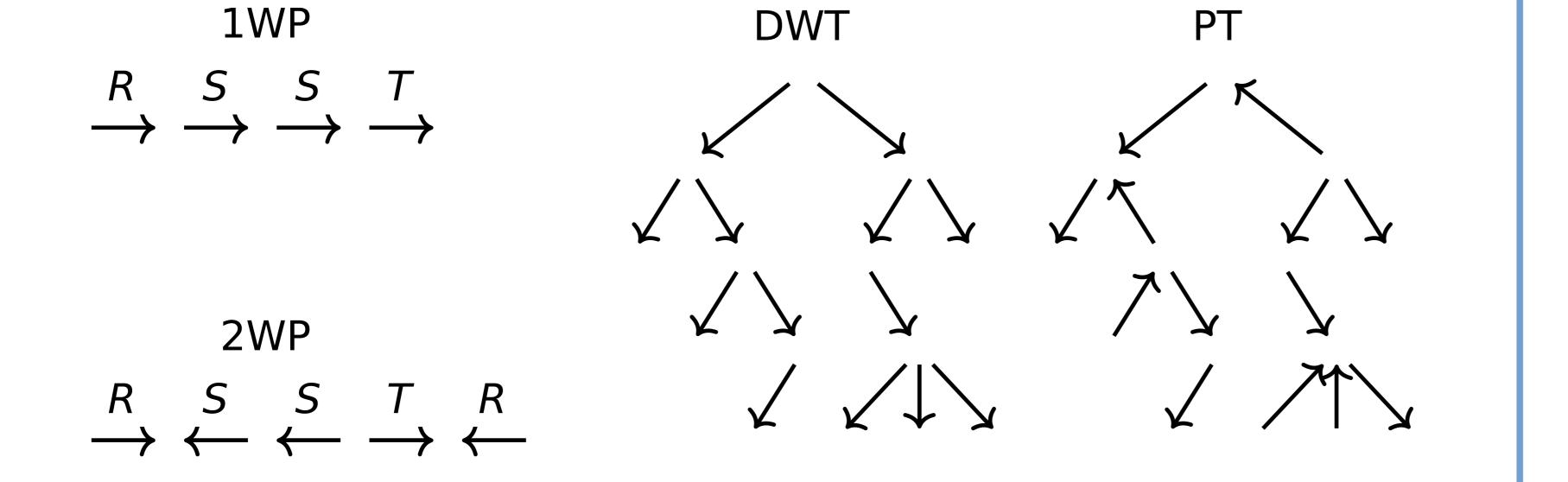
#### **Our Graph Classes**

We intoduce the following graph classes: **one-way paths** (1WP), two-way paths (2WP), downwards trees (DWT) and polytrees (PT)

# **Relationship between classes**







Results	Proof Techniques	Conclusion
Without labels $\downarrow Q  I \rightarrow$ 1WP2WPDWTPTConnected1WP2WPDWTPTConnected2WPPTIME#P-hardPT#P-hard	<ul> <li>Tree automata</li> <li>β-acyclicity</li> <li>X-property</li> <li>Various coding techniques for #P-hardness (#PP2DNF and #Bipartite-Edge-Cover)</li> </ul>	<ul> <li>First study of the combined complexity of Probabilistic Graph Homomorphism</li> <li>Shows the importance of various features</li> <li>Establishes complexity for all combinations of the graph classes we consider</li> </ul>
With >1 labels		However:

$\downarrow Q$	$I \rightarrow$	1WP	2WP	DWT	ΡT	Connected
1V	VP					
2V	VP					
DV	VT		PTIME			
Р	Т					#P-hard
Conn	ected					

References [1] N. Dalvi, D. Suciu The Dichotomy of Probabilistic Inference for Unions of Conjunctive Queries JACM, 2012

[2] A. Amarilli, P. Bourhis, P. Senellart Provenance Circuits for Trees and Treelike Instances Proc. ICALP, 2015

 Graph classes very weak • Nowhere near a dichotomy • Probabilistic equivalent of Feder–Vardi conjecture for combined complexity? • Practical application? (probabilistic databases)